

# Mixing Mutability into the Nanopass Framework

Andy Keep

# Background

- Nanopass framework is a DSL for writing compilers
- Provides a syntax for defining the grammar of an intermediate representation
  - Intermediate representations are immutable\*
  - Mutability can be introduced by adding mutable terminals
  - We will look at using this for variables and basic block labels

\* technically the lists are just Scheme lists, which are mutable

# **A simple compiler**

# Source language

```
(define-language Lsrc
  (terminals
    (datum (d))
    (immediate (imm))
    (symbol (x))
    (primitive (pr)))
  (Expr (e)
    x
    imm
    (quote d)
    (if e0 e1)
    (if e0 e1 e2)
    (and e* ...)
    (or e* ...)
    (not e)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e* ... e)
    (let ([x* e*] ...) e* ... e)
    (letrec ([x* e*] ...) e* ... e)
    (e e* ...)
    (pr e* ...)))
```

# Source language

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(define-language Lsrc
  (terminals
    (datum (d))
    (immediate (imm))
    (symbol (x))
    (primitive (pr)))
  (Expr (e)
    x
    imm
    (quote d)
    (if e0 e1)
    (if e0 e1 e2)
    (and e* ...)
    (or e* ...)
    (not e)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e* ... e)
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    (e e* ...))
    (pr e* ...)))
```

# Source language

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(define-language Lsrc
  (terminals
    (datum (d))
    (symbol (x))
    (primitive (pr)))
  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e* ... e)
    (let ([x* e*] ...) e* ... e)
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    (pr e* ...)))
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  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e* ... e)
    (let ([x* e*] ...) e* ... e)
    (letrec ([x* e*] ...) e* ... e)
    (e e* ...))
  (pr e* ...)))
```

# Source language

```
(define-language Lsrc
  (terminals
    (datum (d))
    (var (x))
    (primitive (pr)))
  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e* ... e)
    (let ([x* e*] ...) e* ... e)
    (letrec ([x* e*] ...) e* ... e)
    (e e* ...))
  (pr e* ...)))
```



# Source language

```
(define-language Lsrc
  (terminals
    (datum (d))
    (var (x))
    (primitive-info (pr)))
  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e* ... e)
    (let ([x* e*] ...) e* ... e)
    (letrec ([x* e*] ...) e* ... e)
    (e e* ...))
    (pr e* ...)))
```

# Source language

```
(define-language Lsrc
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    (datum (d))
    (var (x))
    (primitive-info (pr)))
  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e)
    (let ([x* e*] ...) e)
    (letrec ([x* e*] ...) e)
    (e e* ...)
    (pr e* ...)))
```

# Source language

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    (datum (d))
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    (primitive-info (pr)))
  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e)
    (let ([x* e*] ...) e)
    (letrec ([x* e*] ...) e)
    (e e* ...)
    (pr e* ...)))
```

# Source language

```
(define-language Lsrc
  (terminals
    (datum (d))
    (var (x))
    (primitive-info (pr)))
  (Expr (e)
    x
    (quote d)
    (if e0 e1 e2)
    (set! x e)
    (begin e* ... e)
    (lambda (x* ...) e)
    (let ([x* e*] ...) e)
    (letrec ([x* e*] ...) e)
    (callable e* ...))
  (Callable (callable)
    e
    pr))
```

**Target language?**

# Target language

- LLVM 10
  - A bit lower level than C
  - Better handling of tail calls
  - Brand new (may require installing llvm and clang 10)
  - Required a bit of SSA conversion

# Overall compiler

parse-scheme

convert-complex-datum

uncover-assigned!

purify-letrec

convert-assignments

optimize-direct-call

remove-anonymous-lambda

sanitize-binding-forms

uncover-free

convert-closures

optimize-known-call

introduce-procedure-primitives

lift-letrec

normalize-context

specify-representation

uncover-locals

remove-let

remove-complex-operators

flatten-set!

expose-basic-blocks

optimize-blocks

convert-to-ssa

flatten-functions

eliminate-simple-moves

generate-llvm-code

# Overall compiler

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generate-llvm-code



# Parsing Scheme

- Start with initial environment with syntax and primitives
- Extend environment mapping symbols to a variable record at binding sites
- Replace references to the symbols in the environment with variable records
- Variable records contain a mutable flags field and a mutable "slot"
- References and binding locations share variable record
- No longer need to build environments for variables later
- This is also how Chez Scheme handles variables

# Assignment conversion

# Assignment conversion

```
(let ([x 5] [y 7])  
  (set! x (* x 2))  
  (+ x y))
```

# Assignment conversion

```
(let ([x 5] [y 7])  
  (set! x (* x 2))  
  (+ x y))
```

```
(let ([t 5] [y 7])  
  (let ([x (cons t (void))])  
    (set-car! x (* (car x) 2))  
    (+ (car x) y)))
```

# Assignment conversion

```
(let ([x 5] [y 7])  
  (set! x (* x 2))  
  (+ x y))
```

```
(let ([t 5] [y 7])  
  (let ([x (cons t (void))])  
    (set-car! x (* (car x) 2))  
    (+ (car x) y)))
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# Assignment conversion

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(let ([x 5] [y 7])  
  (set! x (* x 2))  
  (+ x y))
```

```
(let ([t 5] [y 7])  
  (let ([x (cons t (void))])  
    (set-car! x (* (car x) 2))  
    (+ (car x) y)))
```

# Assignment conversion

```
(let ([x 5] [y 7])  
  (set! x (* x 2))  
  (+ x y))  
  
(let ([t 5] [y 7])  
  (let ([x (cons t (void))])  
    (set-car! x (* (car x) 2))  
    (+ (car x) y)))
```

The diagram illustrates the conversion of a mutable variable to a cons cell. On the left, the original code uses `set!` to mutate the variable `x`. On the right, the converted code uses `cons` to create a mutable cell `(car x)` and `set-car!` to mutate it. Blue boxes highlight the corresponding parts: `x` in the first code maps to `(car x)` in the second, and `x` in the second code maps to `(car x)` in the third. Arrows indicate these mappings.

# Assignment conversion

```
(let ([x 5] [y 7])  
  (set! x (* x 2))  
  (+ x y))
```

```
(let ([t 5] [y 7])  
  (let ([x (cons t (void))] )  
    (set-car! x (* (car x) 2))  
    (+ (car x) y)))
```



# Uncover assigned variables

```
(define-pass uncover-assigned! : Ldatum (ir) -> Ldatum ()  
  (Expr : Expr (ir) -> Expr ()  
    [(set! ,x ,e)] (var-flags-assigned-set! x #t) ir]))
```

# Convert assignments

```
(define-pass convert-assignments : Lletrec (ir) -> Lno-assign ()
  (Lambda : Lambda (ir) -> Lambda ()
    [(lambda (,x* ...) ,e)
     (let-values ([(x* e) (convert-bindings x* e)])
       `(lambda (,x* ...) ,e))])
  (Expr : Expr (ir) -> Expr ()
    [,x (if (var-flags-assigned? x) `(,car-pr ,x) x)]
    [(set! ,x ,[e]) `(,set-car!-pr ,x ,e)]
    [(let ([,x* ,[e*]] ...) ,e)
     (let-values ([(x* e) (convert-bindings x* e)])
       `(let ([,x* ,e*] ...) ,e))])])
```

# Convert assignments

```
(define convert-bindings
  (lambda (x* e)
    (with-assigned x*
      (case-lambda
        [(x*) (values x* (Expr e))]
        [(x* assigned-x* new-x*)
         (values x*
                  (with-output-language (Lno-assign Expr)
                    (let ([pr* (map
                              (lambda (new-x)
                                `(,cons-pr ,new-x (,void-pr)))
                              new-x*))
                      `(let ([,assigned-x* ,pr*] ...)
                        ,(Expr e))))))]))))
```

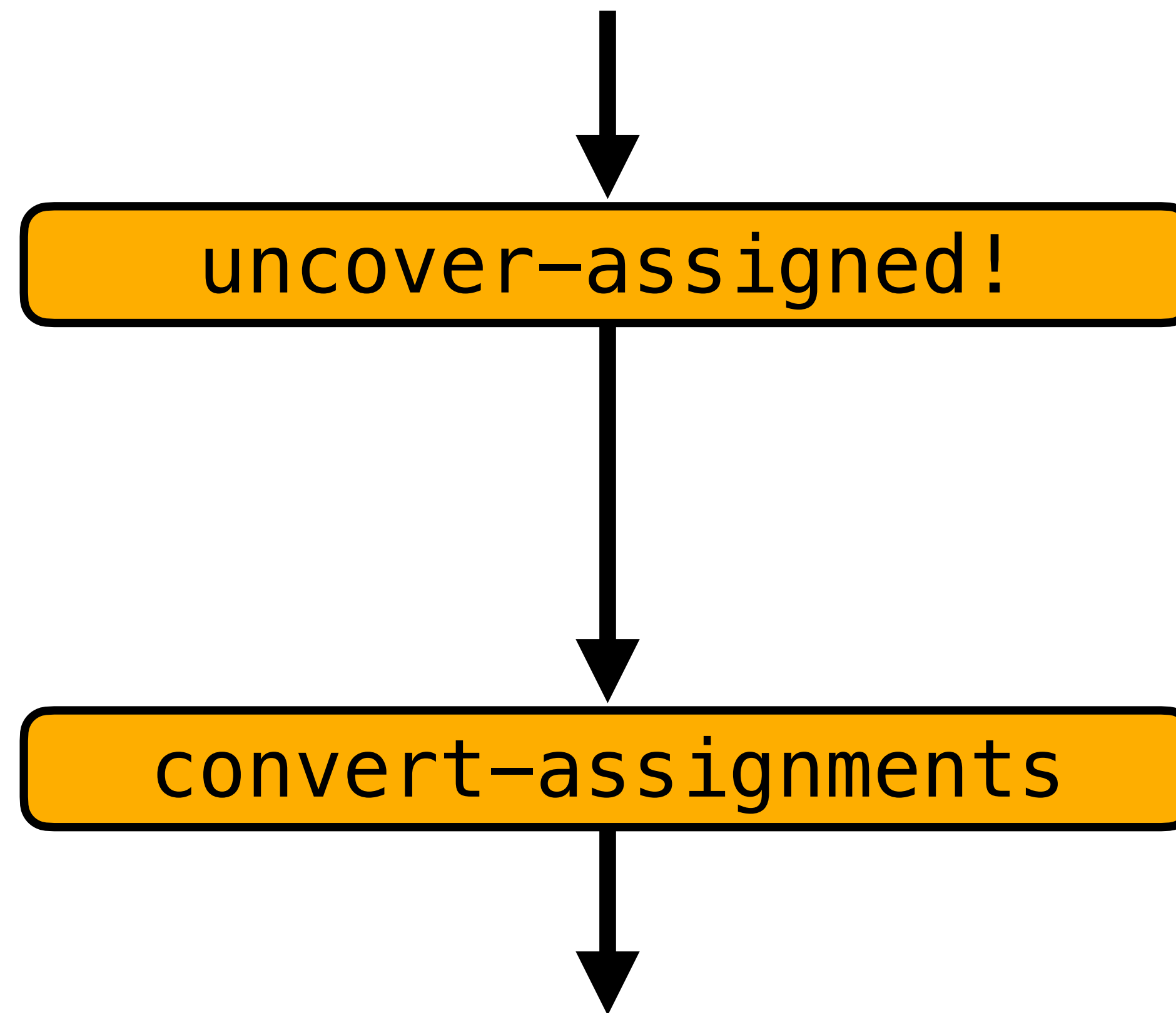
# Convert assignments

```
(define with-assigned
  (lambda (x* f)
    (let l ([x* x*] [rx* '()] [rset-x* '()] [rnew-x* '()])
      (if (null? x*)
          (if (null? rset-x*)
              (f (reverse rx*))
              (f (reverse rx*) (reverse rset-x*)
                 (reverse rnew-x*)))
          (let ([x (car x*)] [x* (cdr x*)])
            (if (var-flags-assigned? x)
                (let ([new-x (make-var x)])
                  (l x* (cons new-x rx*)
                     (cons x rset-x*) (cons new-x rnew-x*)))
                (l x* (cons x rx*) rset-x* rnew-x*)))))))
```

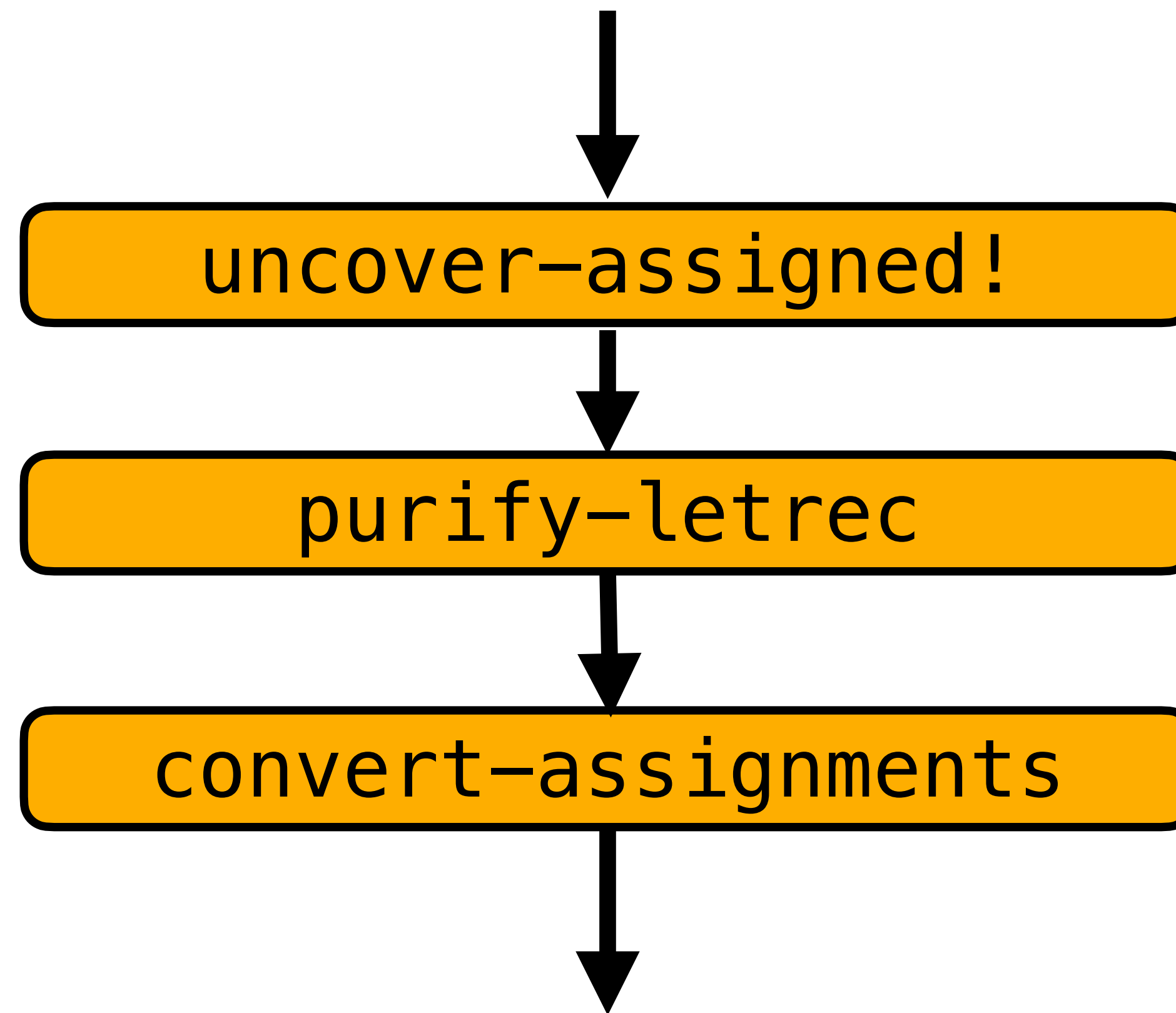
# Convert assignments

```
(define-pass convert-assignments : Lletrec (ir) -> Lno-assign ()
  (Lambda : Lambda (ir) -> Lambda ()
    [(lambda (,x* ...) ,e)
     (let-values ([(x* e) (convert-bindings x* e)])
       `(lambda (,x* ...) ,e))]))
  (Expr : Expr (ir) -> Expr ()
    [,x (if (var-flags-assigned? x) `(,car-pr ,x) x)]
    [(set! ,x ,[e]) `(,set-car!-pr ,x ,e)]
    [(let ([,x* ,[e*]] ...) ,e)
     (let-values ([(x* e) (convert-bindings x* e)])
       `(let ([,x* ,e*] ...) ,e))]))
```

# One small problem



# One small problem



# Purify letrec

- Categorizes **letrec** bindings into: assigned, simple, lambda, and complex
- Assigned are already marked assigned, no problem there
- Simple and lambda are not assigned, and don't become assigned
- Complex on the other hand, become assigned where they were not before
- We need to track this assignment.



# Purify letrec

```
(cond
  [(var-flags-assigned? x)
   (loop (cdr tx*) (cdr e*) xs* es* xl* el*
         (cons x xc*) (cons (Expr e) ec*))]
  [(lambda-expr? e)
   (loop (cdr tx*) (cdr e*) xs* es* (cons x xl*)
         (cons (Expr e) el*) xc* ec*)]
  [(simple-expr? e)
   (loop (cdr tx*) (cdr e*) (cons x xs*)
         (cons (Expr e) es*) xl* el* xc* ec*)]
  [else
   ;; we made an unassigned variable assigned, mark it.
   (var-flags-assigned-set! x #t)
   (loop (cdr tx*) (cdr e*) xs* es* xl* el*
         (cons x xc*) (cons (Expr e) ec*))])])
```

# Purify letrec

```
(cond
  [(var-flags-assigned? x)
   (loop (cdr tx*) (cdr e*) xs* es* xl* el*
         (cons x xc*) (cons (Expr e) ec*))])
 [(lambda-expr? e)
  (loop (cdr tx*) (cdr e*) xs* es* (cons x xl*)
        (cons (Expr e) el*) xc* ec*))])
 [(simple-expr? e)
  (loop (cdr tx*) (cdr e*) (cons x xs*)
        (cons (Expr e) es*) xl* el* xc* ec*))])
 [else
  ;; we made an unassigned variable assigned, mark it.
  (var-flags-assigned-set! x #t)
  (loop (cdr tx*) (cdr e*) xs* es* xl* el*
        (cons x xc*) (cons (Expr e) ec*))])])
```

# Closure conversion

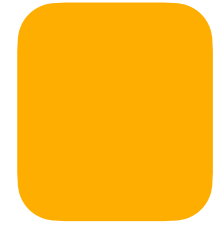


# Free variable analysis

```
(lambda (x)
  (lambda (y)
    (lambda (z)
      (+ x (+ y z))))))
```

# Free variable analysis

```
(lambda (x)  
  (lambda (y) x
    (lambda (z) x y
      (+ x (+ y z))))))
```

# Free variable analysis

( lambda ( x<sup>0</sup> )   
 ( lambda ( y<sup>1</sup> )  x  
 ( lambda ( z<sup>2</sup> )  x y  
 ( + x<sup>0</sup> ( + y<sup>1</sup> z<sup>2</sup> ) ) ) ) ) )

# Free variable analysis

( lambda ( x<sup>0</sup> )   0 0 0  
  ( lambda ( y<sup>1</sup> ) x 0 0 1  
    ( lambda ( z<sup>2</sup> ) x y 0 1 1  
      ( + x<sup>0</sup> ( + y<sup>1</sup> z<sup>2</sup> ) ) ) ) ) )

# Free variable analysis

( lambda ( x<sup>0</sup> )   0 0 0 |  
  ( lambda ( y<sup>1</sup> ) x 0 0 | 1  
    ( lambda ( z<sup>2</sup> ) x y 0 | 1 1  
      ( + x<sup>0</sup> ( + y<sup>1</sup> z<sup>2</sup> ) ) ) ) ) )



# Uncover free

```
(define-pass uncover-free : Lsanitized (ir) -> Lfree ()
  (Callable : Callable (e index fv-info) -> Callable ())
  (Expr : Expr (e index fv-info) -> Expr ())
  [,x (record-ref! x fv-info) x]
  [(let ([,x* ,[e*]] ...) ,e)
   (with-offsets (index x*)
    `(let ([,x* ,e*] ...) ,(Expr e index fv-info)))]
  [(letrec ([,x* ,f*] ...) ,e)
   (with-offsets (index x*)
    (let ([f* (map (lambda (f) (Lambda f index fv-info)) f*)]
          [e (Expr e index fv-info)])
      `(letrec ([,x* ,f*] ...) ,e)))]
  (Lambda : Lambda (e index outer-fv-info) -> Lambda ())
  [(lambda (,x* ...) ,e)
   (let ([fv-info (make-fv-info index)])
     (with-offsets (index x*)
      (let ([e (Expr e index fv-info)])
        (let ([fv* (fv-info-fv* fv-info)]
              (for-each (lambda (fv) (record-ref! fv outer-fv-info)) fv*))
          `(lambda (,x* ...) (free (,fv* ...) ,e)))))))]
  (Expr ir 0 (make-fv-info 0)))
```

# Uncover free

```
(define-pass uncover-free : Lsanitized (ir) -> Lfree ())
(Ca
(Ex (define (set-offsets! x* index)
  [ (fold-left
    [ (lambda (index x)
      [ (var-slot-set! x index)
        [ (fx+ index 1))
          index x*))
    (define ($with-offsets index x* p)
      (let ([index (set-offsets! x* index)])
        (let ([v (p index)])
          (for-each (lambda (x) (var-slot-set! x #f)) x*)
            v)))
    (La
  [ (define-syntax with-offsets
    [ (lambda (x)
      [ (syntax-case x ()
        [ (_ (index ?x*) ?e ?es ...)
          [ (identifier? #'index)
            #'($with-offsets index ?x* (lambda (index) ?e ?es ...))]]))
    (Expr ir 0 (make-fv-info 0)))
```

# Uncover free

```
(define-pass uncover-free : Lsanitized (ir) -> Lfree ()
  (Callable : Callable (e index fv-info) -> Callable ()))
(Expr
  [ (define-record-type fv-info
    [ (nongenerative)
      (fields lid (mutable mask) (mutable fv*))
      (protocol
        [ (lambda (new)
            (lambda (index)
              (new index 0 '()))))
          (define (record-ref! x info)
            (let ([idx (var-slot x)])
              (when (fx<? idx (fv-info-lid info))
                (let ([mask (fv-info-mask info)])
                  (unless (bitwise-bit-set? mask idx)
                    (fv-info-mask-set! info (bitwise-copy-bit mask idx 1))
                    (fv-info-fv*-set! info (cons x (fv-info-fv* info))))))))
            (lambda (,x* ...) (free (,fv* ...) ,e))))))
    (Expr ir 0 (make-fv-info 0)))
```

# Uncover free

```
(define-pass uncover-free : Lsanitized (ir) -> Lfree ()
  (Callable : Callable (e index fv-info) -> Callable ())
  (Expr : Expr (e index fv-info) -> Expr ())
  [,x (record-ref! x fv-info) x]
  [(let ([,x* , [e*]] ...) ,e)
   (with-offsets (index x*)
    `(let ([,x* ,e*] ...) ,(Expr e index fv-info)))]
  [(letrec ([,x* ,f*] ...) ,e)
   (with-offsets (index x*)
    (let ([f* (map (lambda (f) (Lambda f index fv-info)) f*)]
          [e (Expr e index fv-info)])
      `(letrec ([,x* ,f*] ...) ,e)))]
  (Lambda : Lambda (e index outer-fv-info) -> Lambda ())
  [(lambda (,x* ...) ,e)
   (let ([fv-info (make-fv-info index)])
     (with-offsets (index x*)
      (let ([e (Expr e index fv-info)])
        (let ([fv* (fv-info-fv* fv-info)]
              (for-each (lambda (fv) (record-ref! fv outer-fv-info)) fv*))
          `(lambda (,x* ...) (free (,fv* ...) ,e)))))))]
  (Expr ir 0 (make-fv-info 0)))
```

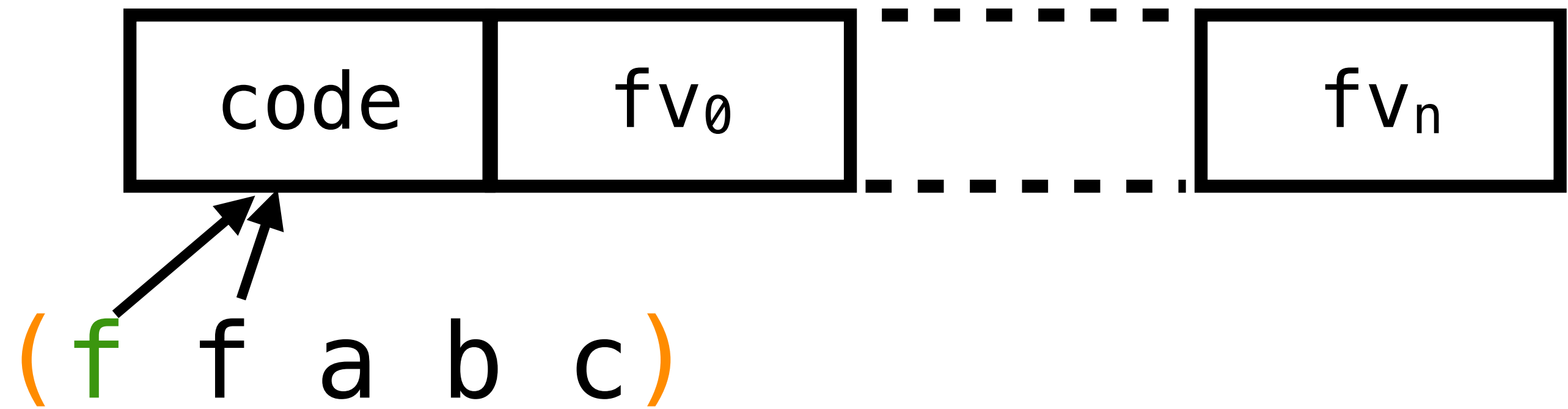
# Compiling function calls

( f a b c )

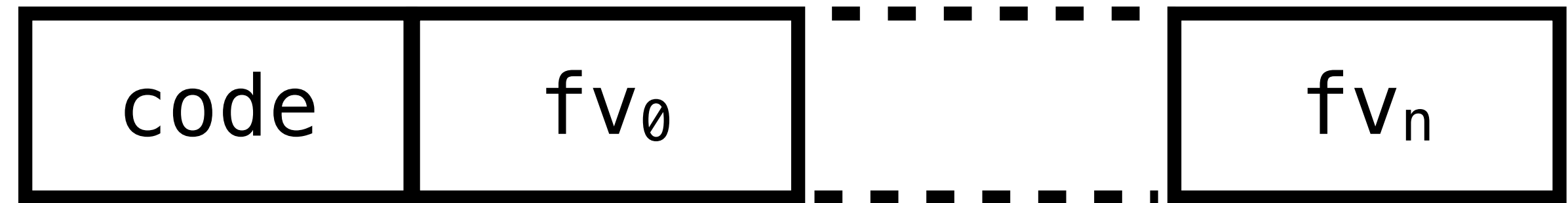
# Compiling function calls

( f f a b c )

# Compiling function calls



# Compiling function calls



`(( $procedure-code f ) f a b c )`

Two arrows point from the `f` tokens in the code above to the `fv0` and `fvn` boxes in the diagram above.



# Compiling function calls

```
(letrec ([lf (lambda (x y z) ---)])  
  (closures ([f lf ---])  
            ---  
            (lf f a b c)  
            ---)))
```

# Optimize known call

```
(define-pass optimize-known-call : Lclosure (ir) -> Lclosure ()
  (Lambda : Lambda (f) -> Lambda ())
  (Expr : Expr (ir) -> Expr ())
  [( ,x , [e*] ... )
    (cond
      [(var-slot x) => (lambda (l) `( ,l ,e* ...))]
      [else `( ,x ,e* ...)])]
  [(letrec ([ ,l0* , f*] ... )
    (closures ([ ,x* , l* , x** ... ] ... ) ,e))
    (for-each (lambda (x l) (var-slot-set! x l)) x* l*)
    (let ([f* (map Lambda f*)] [e (Expr e)])
      (for-each (lambda (x) (var-slot-set! x #f)) x*)
      `(letrec ([ ,l0* , f*] ... )
        (closures ([ ,x* , l* , x** ... ] ... ) ,e)))]
  ;; NB: should be unnecessary
  [(letrec ([ ,l* , f*] ... ) ,clbody) (errorf who "unreachable"])]])
```

# Introduce procedure primitives

```
(letrec ([lf (lambda (cp z)
             (bind-free (cp x y)
                        (+ x (+ y z))))])
 (closures ([f lf x y]
            f)))
```

# Introduce procedure primitives

```
(letrec ([lf (lambda (cp z)
            (bind-free (cp x y)
                        (+ x (+ y z))))])
  (closures ([f lf x y]
            f)))
```

```
(letrec ([lf (lambda (cp z)
            (+ ($procedure-ref cp '0)
              (+ ($procedure-ref cp '1)
                z))))])
  (let ([f ($make-closure lf '2)])
    ($procedure-set! f '0 x)
    ($procedure-set! f '1 y)
    f)))
```

# Introduce procedure primitives

```
(letrec ([lf (lambda (cp z)
            (bind-free (cp x y)
                       (+ x (+ y z))))])
  (closures ([f lf x y])
            f))
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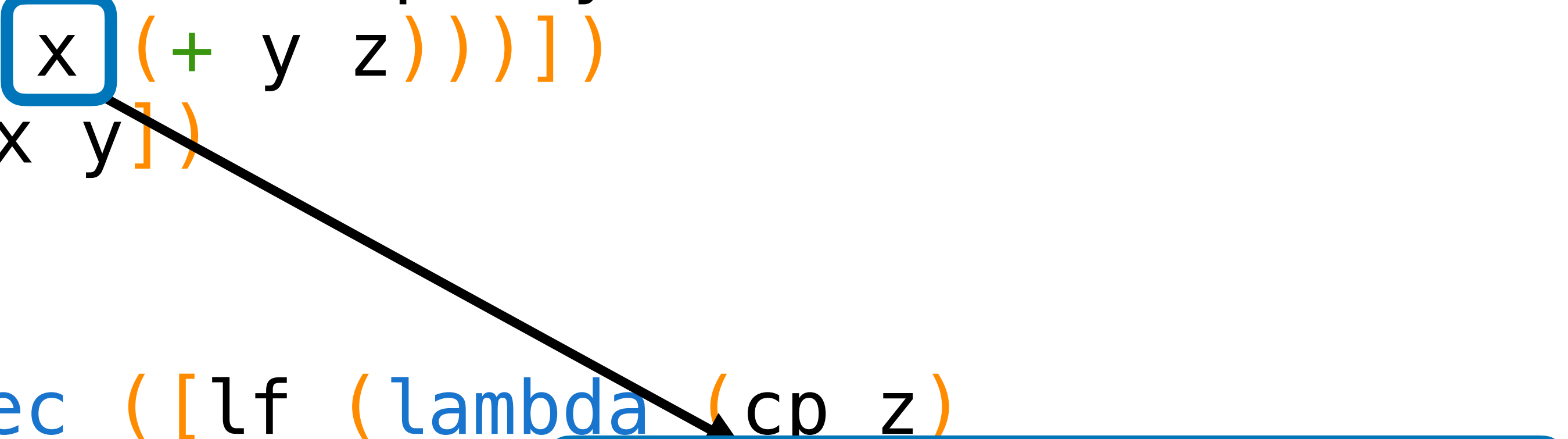
```
(letrec ([lf (lambda (cp z)
            (+ ($procedure-ref cp '0)
              (+ ($procedure-ref cp '1)
                z))))])
```

```
(let ([f ($make-closure lf '2)])
  ($procedure-set! f '0 x)
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```

# Introduce procedure primitives

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(letrec ([lf (lambda (cp z)
            (bind-free (cp x y)
                        (+ x (+ y z))))])
  (closures ([f lf x y]
            f)))
```

```
(letrec ([lf (lambda (cp z)
            (+ ($procedure-ref cp '0)
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  (let ([f ($make-closure lf '2)])
    ($procedure-set! f '0 x)
    ($procedure-set! f '1 y)
    f)))
```



# Introduce procedure primitives

```
(letrec ([lf (lambda (cp z)
            (bind-free (cp x y)
                        (+ x (+ y z))))])
  (closures ([f lf x y]
            f)))
```

```
(letrec ([lf (lambda (cp z)
            (+ ($procedure-ref cp '0)
              (+ ($procedure-ref cp '1)
                z))))])
```

```
(let ([f ($make-closure lf '2)])
  ($procedure-set! f '0 x)
  ($procedure-set! f '1 y)
  f))
```

# Introduce procedure primitives

```
(define-pass introduce-procedure-primitives : Lclosure (ir) -> Lproc ()
  (var : var (x) -> Expr ())
  (cond
    [(var-slot x) => build-procedure-ref]
    [else x]))
(Expr : Expr (e) -> Expr ())
[,x (var x)]
[(letrec ([,l0* ,f*] ...)
  (closures ([,x* ,l1* ,e**] ...) ,e))]
` (letrec ([,l0* ,f*] ...)
  (let ([,x* ,(build-make-proc! l1* e**)] ...)
    ,(build-procedure-set! x* e** e)))
[(,l ,e* ...) `(,l ,e* ...)]
[(,pr ,e* ...) `(,pr ,e* ...)]
[(,[e] ,e* ...) `((,procedure-code-pr ,e) ,e* ...)])
(Lambda : Lambda (f) -> Lambda ())
[(lambda (,x* ...) (bind-free (,x ,x0* ...) ,e))
 (with-fv* x x0* (lambda () `(lambda (,x* ...) ,(Expr e)))))]
```



# Introduce procedure primitives

```
(define (build-procedure-ref pr)
  `(,procedure-ref-pr ,(car pr) (quote ,(cdr pr))))
(va
  (define (build-make-proc! l* e**)
    (map
      (lambda (l e*)
        `(,make-procedure-pr ,l (quote ,(length e*))))
      l* e**))
(Ex)
[
  (define (build-procedure-set! x* e** e)
    (let ([ps* (fold-right
              (lambda (x e* ps*)
                (fold-right
                  (lambda (e i ps*)
                    (cons `(,procedure-set!-pr ,x (quote ,i) ,e) ps*))
                  ps* e* (enumerate e*))
                '() x* e**)]])
      (if (null? ps*)
          e
          `(begin ,ps* ... ,e))))
(La)
[
```

# Introduce procedure primitives

```

(define-pass introduce-procedure-primitives : Lclosure (ir) -> Lproc ()
  (var : var (x) -> Expr ())
  (cond
    [(var-slot x
      [else x]))
    (Expr : Expr (e)
      [, x (var x)]
      [(letrec ([, l0
        (closures (
          ` (letrec ([, l
            (let ([, x*
              ,(build-
                [, l , [e*] ...
                [, pr , [e*] ..
                [, [e] , [e*] .
        (Lambda : Lambda
          [(lambda (, x* ...) (bind-free (, x , x0* ...) , e))]
          (with-fv* x x0* (lambda () `(lambda (, x* ...) ,(Expr e)))))]))

```

**Optimize and reorder blocks**

# Optimize blocks

( labels ( [ , l\* , t\* ] ... ) , l )

# Optimize blocks

```
(labels ([, l* , t*] ...), l)
```

```
(build-graph! l* t*)
```

# Optimize blocks

```
(labels ([, l* , t*] ... ) , l)  
  
(for-each  
  (lambda (l t)  
    (label-slot-set! l  
      (make-graph-node t)))  
  l* t*)
```

# Optimize blocks

```
(let loop ([wl (list l)] [rl* '()] [rt* '()])
  (if (null? wl)
      (begin
        (for-each (lambda (l) (label-slot-set! l #f)) l*)
        `(labels ([, (reverse rl*) , (reverse rt*)] ...), l))
      (let ([l (car wl)] [wl (cdr wl)])
        (let ([node (label-slot l)])
          (if (graph-node-written? node)
              (loop wl rl* rt*)
              (begin
                (graph-node-written?-set! node #t)
                (let-values ([ (t wl)
                              (rewrite-tail
                               (graph-node-tail node) wl)])
                  (loop wl (cons l rl*) (cons t rt*))))))))))
```

# Optimize blocks

```
(rewrite-tail : Tail (t wl) -> Tail (wl))
(let l
  (if
    [(begin ,ef* ... ,t)
     (let*-values ([ (ef* wl) (rewrite-effect* ef* wl)]
                  [(t wl) (rewrite-tail t wl)])
       (values `(begin ,ef* ... ,t) wl)))
     [(goto ,l)
      (let ([l (extract-final-target l)])
        (values `(goto ,l) (extend-worklist l wl)))]
     [(return ,l)
      (let ([l (extract-final-target l)])
        (values `(return ,l) (extend-worklist l wl)))]
     [(return ,tr) (values `(return ,tr) wl)]
     [(if (,relop ,tr0 ,tr1) (,l0) (,l1))
      (let ([l0 (extract-final-target l0)]
            [l1 (extract-final-target l1)])
        (values `(if (,relop ,tr0 ,tr1) (,l0) (,l1))
                 (extend-worklist l0 l1 wl)))]))])
```



# Other uses of mutation

# Mutation in the compiler

- **convert-complex-datum** uses **fluid-let** for creating constant bindings
- **lift-letrec** use **fluid-let** for binding top-level labels and functions
- **uncover-locals** uses **fluid-let** for binding locals list
- **remove-complex-opera\*** uses **fluid-let** for binding locals list
- **expose-basic-blocks** uses **fluid-let** for binding locals list

# Mutation in the compiler

- **convert-to-ssa** uses var slot for variable renaming
- **convert-to-ssa** uses multiply assigned flag to find variables that need phi
- **convert-to-ssa** use label slot for creating control-flow graph
- **eliminate-simple-moves** uses var slot for replacing reference with value

**Wrapping up**

# Wrapping up

- Limited and controlled use of mutable storage can be useful
- Mutable information that lasts across passes needs to be maintained
- When using a mutable cell for a single pass, we must cleanup at the end
- We assume only one thread will have a program at a given time
- We can avoid the cost of reconstructing environments using records
- You can try it out yourself:  
<https://github.com/akeep/scheme-to-llvm>

# Thanks!

<https://github.com/akeep/scheme-to-llvm>

# Questions?

<https://github.com/akeep/scheme-to-llvm>